

Matching Theory and Practice

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Overview

1. Introduction
2. Matching Theory
 - Marriage Market
 - Stability
 - Deferred-Acceptance
3. Applications
 - School Choice
 - Kindergarten in Victoria
4. My Research
 - Matching with Quantity
 - Refugee Dispersal
5. Conclusion

Matching Markets

- Money is extremely useful to facilitate transactions
 - ▶ Price equilibrates supply and demand
 - ▶ Markets organise themselves well
 - ▶ Adam Smith's invisible hand
- In matching markets, money cannot be used
 - ▶ There may be a price but it does not equilibrate supply and demand
 - ▶ These markets do not perform well if left to themselves (market failure)
 - ▶ Economists can redesign these markets to make them work better
- Examples
 - ▶ School or university admission
 - ▶ Kidney donations
 - ▶ Allocations of tasks within an organisation
 - ▶ Refugee resettlement

A Brief History of Matching

- Gale and Shapley (1962)
 - ▶ Brilliant and easy to read paper
 - ▶ Theoretical exercise about an abstract marriage market

- Real world applications
 - ▶ Started in early 2000's
 - ▶ Very active field since then

- 2012 Nobel Prize in Economics
 - ▶ Lloyd Shapley and Al Roth
 - ▶ “Who Gets What and Why?”

Marriage Market (GS 1962)

- Set of women $\{w_1, w_2, \dots, w_n\}$ and set of men $\{m_1, m_2, \dots, m_n\}$
 - ▶ Each woman can be matched (married) to at most one man
 - ▶ Each man can be matched (married) to at most one woman
- People care who they marry
 - ▶ Women have (ordinal) preferences over men and remaining single
 - ▶ Men have (ordinal) preferences over women and remaining single
- How do we best match these men and women?
 - ▶ A key concept is stability
 - ▶ It ensures people do not want to rematch
 - ▶ Essential to the success of two-sided matching markets

Stability

Definition (Individual Rationality, GS 62)

A matching is **individually rational** if there does not exist any woman or man who would prefer to remain single than to be matched with his/her current partner.

Definition (Stability, GS 62)

A matching is **stable** if it is individually rational there does not exist any woman and any man who would both prefer to be matched with each other than with their current partners.

Example

- Consider the following matching:

$$(w_1, m_1), (w_2, m_2)$$

- **Individual rationality** requires

- ▶ w_1 prefers to be with m_1 than single
- ▶ m_1 prefers to be with w_1 than single
- ▶ w_2 prefers to be with m_2 than single
- ▶ m_2 prefers to be with w_2 than single

- **Stability** requires

- ▶ Individual rationality
- ▶ EITHER w_1 prefers m_1 to m_2 OR m_2 prefers w_2 to w_1
- ▶ EITHER w_2 prefers m_2 to m_1 OR m_1 prefers w_1 to w_2

Deferred-Acceptance Algorithm

- Each man proposes to the woman he prefers (if any)
 - ▶ Each woman tentatively accepts her favourite proposal (if any)
 - ▶ She rejects all other proposals
- Each man makes a new proposal
 - ▶ If he was accepted he proposes to the same woman again
 - ▶ If he was rejected he proposes to his next favourite woman (if any)
- The algorithm terminates when all proposals are accepted
 - ▶ Each man is matched with the woman to whom he last proposed
 - ▶ Each man who did not make a proposal remains single
 - ▶ Each woman who did not accept any proposal remains single
- The algorithm is simple and easy to use in practice
 - ▶ It can be coded in an Excel spreadsheet (Visual Basics)

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 1

Eric \rightarrow Ashley
 Frank \rightarrow Ashley
 George \rightarrow Barbara
 Henry \rightarrow Chelsea

In the first round, each man proposes to his favourite woman.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 1

Eric \rightarrow Ashley ✓
 Frank \rightarrow Ashley ✗
 George \rightarrow Barbara
 Henry \rightarrow Chelsea

Ashley chooses Eric over Frank.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 1

Eric \rightarrow Ashley ✓
 Frank \rightarrow Ashley ✗
 George \rightarrow Barbara ✗
 Henry \rightarrow Chelsea

Barbara rejects George's proposal.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 1

Eric \rightarrow Ashley ✓
 Frank \rightarrow Ashley ✗
 George \rightarrow Barbara ✗
 Henry \rightarrow Chelsea ✓

Chelsea tentatively accepts Harry's proposal.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 1

Eric	→	Ashley	✓
Frank	→	Ashley	✗
George	→	Barbara	✗
Henry	→	Chelsea	✓

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 1

Eric	→	Ashley	✓
Frank	→	Ashley	✗
George	→	Barbara	✗
Henry	→	Chelsea	✓

In Round 2, Frank and George will propose to Chelsea.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 2

Eric \rightarrow Ashley
 Frank \rightarrow Chelsea
 George \rightarrow Chelsea
 Henry \rightarrow Chelsea

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 2

Eric \rightarrow Ashley ✓
 Frank \rightarrow Chelsea
 George \rightarrow Chelsea
 Henry \rightarrow Chelsea

Ashley again tentatively accepts Eric's proposal.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 2

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Chelsea	✗
Henry	→	Chelsea	✗

Chelsea chooses Frank over George and Henry.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 2

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Chelsea	✗
Henry	→	Chelsea	✗

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 2

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Chelsea	✗
Henry	→	Chelsea	✗

In Round 3, George and Henry will propose to Dory.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 3

Eric \rightarrow Ashley
 Frank \rightarrow Chelsea
 George \rightarrow Dory
 Henry \rightarrow Dory

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 3

Eric \rightarrow Ashley ✓
 Frank \rightarrow Chelsea
 George \rightarrow Dory
 Henry \rightarrow Dory

Ashley tentatively accept Eric's proposal.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 3

Eric → Ashley ✓
 Frank → Chelsea ✓
 George → Dory
 Henry → Dory

Chelsea tentatively accept Frank's proposal.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 3

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Dory	✓
Henry	→	Dory	✗

Dory chooses George over Henry.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 3

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Dory	✓
Henry	→	Dory	✗

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 3

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Dory	✓
Henry	→	Dory	✗

In Round 4, Henry will propose to Ashley.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 4

Eric \rightarrow Ashley
 Frank \rightarrow Chelsea
 George \rightarrow Dory
 Henry \rightarrow Ashley

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 4

Eric \rightarrow Ashley ✓
 Frank \rightarrow Chelsea
 George \rightarrow Dory
 Henry \rightarrow Ashley ✗

Ashley chooses Eric over Henry.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 4

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Dory	✓
Henry	→	Ashley	✗

Chelsea and Dory tentatively accept their respective proposals.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 4

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Dory	✓
Henry	→	Ashley	✗

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 4

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Dory	✓
Henry	→	Ashley	✗

Henry has run out of options and will not make any proposal in Round 5.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 5

Eric \rightarrow Ashley
 Frank \rightarrow Chelsea
 George \rightarrow Dory
 Henry \rightarrow \emptyset

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Round 5

Eric	→	Ashley	✓
Frank	→	Chelsea	✓
George	→	Dory	✓
Henry	→	\emptyset	✓

All proposals are accepted and the algorithm terminates.

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Outcome

Eric is matched with Ashley

Frank is matched with Chelsea

George is matched with Dory

Henry and Barbara remain unmatched

Example

Ashley:	G, E, H, F, \emptyset	Eric:	A, \emptyset
Barbara:	E, H, \emptyset	Frank:	A, C, B, D, \emptyset
Chelsea:	F, H, G, \emptyset	George:	B, C, D, A, \emptyset
Dory:	F, G, H, E, \emptyset	Henry:	C, D, A, \emptyset

Summary

Round 1	Round 2	Round 3	Round 4	Round 5
$E \rightarrow A \checkmark$	$E \rightarrow A \checkmark$	$E \rightarrow A \checkmark$	$E \rightarrow A \checkmark$	$E \rightarrow A \checkmark$
$F \rightarrow A \times$	$F \rightarrow C \checkmark$	$F \rightarrow C \checkmark$	$F \rightarrow C \checkmark$	$F \rightarrow C \checkmark$
$G \rightarrow B \times$	$G \rightarrow C \times$	$G \rightarrow D \checkmark$	$G \rightarrow D \checkmark$	$G \rightarrow D \checkmark$
$H \rightarrow C \checkmark$	$H \rightarrow C \times$	$H \rightarrow D \times$	$H \rightarrow A \times$	$H \rightarrow \emptyset \checkmark$

Properties of DA

Theorem

The matching produced by the man-proposing deferred-acceptance algorithm is the man-optimal stable matching.

- Man-optimal stable matching
 - ▶ In any other stable matching, all men are either matched with the same woman or with one they like less
 - ▶ Best stable matching from the men's point of view
 - ▶ Worst stable matching from the women's point of view

Properties of DA

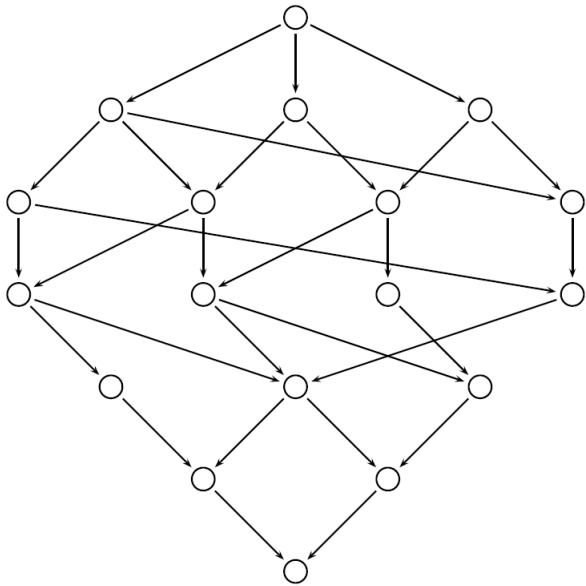
Theorem

The matching produced by the woman-proposing deferred-acceptance algorithm is the woman-optimal stable matching.

- Woman-optimal stable matching
 - ▶ In any other stable matching, all women are either matched with the same man or with one they like less
 - ▶ Best stable matching from the women's point of view
 - ▶ Worst stable matching from the men's point of view

Set of Stable Matchings

- On one extreme, men-optimal stable matching.
 - ▶ Found by the men-proposing DA
 - ▶ Best stable matching for men, worst for women
- On the other extreme, woman-optimal stable matching
 - ▶ Found by the women-proposing DA
 - ▶ Best stable matching for women, worst for men
- The set of stable matching is always nonempty
 - ▶ If both versions of DA give the same matching: unique stable matching
 - ▶ Otherwise DA gives the two extremes
 - ▶ There may be more stable matchings in between



Incentive Properties

Theorem (GS 62)

*The deferred acceptance is **strategy-proof** for the **proposing** side but not for the proposed side.*

- In the men-proposing deferred-acceptance algorithm:
 - ▶ Men can only lose out if they misrepresent their preferences
 - ▶ Women can potentially gain by misrepresenting their preferences
- Finding the right strategy is difficult and risky
 - ▶ More likely to lose than gain
 - ▶ Generally not regarded as a big problem

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From Theory To Practice

- Gale and Shapley (1962) proposed a theoretical model
 - ▶ To the best of my knowledge no marriage is arranged in this way
 - ▶ Mathematics and romance do not always get along...
 - ▶ The literature remained essentially theoretical until the early 2000's
- Since then many applications
 - ▶ **School Choice**
 - ▶ Kidney Exchange
 - ▶ National Resident Matching Program
- This presentation focuses on school choice
 - ▶ It constitutes the starting point of the applied matching literature
 - ▶ The problem is similar to the marriage market
 - ▶ It is relevant to Victoria

School Choice

- Abdulkadiroglu and Sonmez (2003)
 - ▶ Excellent paper, easy to read
 - ▶ High school students assigned to schools in Boston
 - ▶ It has been extended to many US cities
 - ▶ It could be applied to Melbourne
- Assigning students to their neighbourhood causes problems
 - ▶ Wealthy parents move to areas with good schools
 - ▶ United States cities are very segregated
- Allowing students to choose has three advantages
 - ▶ It is welfare enhancing
 - ▶ It reduces the importance of family wealth
 - ▶ It mixes populations

The Problem

- School Choice existed but was not done optimally
 - ▶ Ineffective “Boston” algorithm
 - ▶ Incentive problem and unfair matching
- The authors proposed a new design
 - ▶ Based on Gale and Shapley (1962)
 - ▶ Uses the deferred acceptance algorithm
- The new design was implemented
 - ▶ Economists have been designing matching markets ever since

The Model

- Very close to the marriage market
 - ▶ Set of students and set of schools
 - ▶ Students have ordinal preferences over schools
 - ▶ Schools have ordinal *priorities* over students
- Many-to-one matching
 - ▶ Each school is matched with many students
 - ▶ Each school has a capacity limit (number of students it can fit)
 - ▶ This hardly makes a difference, GS 62 considered it as an extension
- The market is one-sided
 - ▶ Schools are not strategic agents, school seats are goods
 - ▶ Only students' welfare matters
 - ▶ Schools priorities (not preferences) are determined by law
 - ▶ This is important

One-Sided Market

- Priorities determined by law
 - ▶ Higher priority if the school is in the same neighbourhood
 - ▶ Higher priority if the sibling is attending the school
 - ▶ Lottery
- Stability means fairness
 - ▶ Schools are not strategic agent, they will not rematch
 - ▶ A stable matching is fair: if a student misses out on a school (s)he likes, then all students attending that school have a higher priority
- Student proposing DA has desirable properties
 - ▶ It is strategy-proof
 - ▶ It is stable (fair)
 - ▶ It maximises welfare given the stability (fairness) constraint

Policy Implications

- In the United States
 - ▶ “Boston” algorithm was replaced by deferred acceptance
 - ▶ Similar designs were implemented in other cities

- Can we learn from this in Victoria?
 - ▶ Kindergarten
 - ▶ Schools?
 - ▶ Child Care?

Kindergarten in Victoria

- What is kindergarten?
 - ▶ Often called *Preschool*
 - ▶ One year program, two years before Grade 1
 - ▶ Attendance is optional and places are not guaranteed
 - ▶ Funded by the state, often owned and operated by councils
 - ▶ Sometimes privately owned but strictly regulated
- A matching market
 - ▶ Children (or their parents) have preferences over kindergarten
 - ▶ Priorities for each kindergarten are determined by law
 - ▶ Each kindergarten has a capacity limit
 - ▶ The problem is almost identical to school choice

Using Matching Theory

- Typical process
 - ▶ Centralised at the council level
 - ▶ Four rounds of offers over two months
 - ▶ Outcome is similar to the “Boston” algorithm...
 - ▶ But it takes two months instead of thirty seconds
- It could be replaced by the deferred acceptance algorithm
 - ▶ Large amount of time and paperwork saved
 - ▶ Better allocation
 - ▶ Strategy-proof for families
 - ▶ Better information on demand for kindergarten

Obstacles

- People like to be in control
 - ▶ They are rightfully weary of mysterious algorithms
 - ▶ Explaining how it works goes a long way
- Councils may feel power is taken away from them
 - ▶ They retain control over priorities
 - ▶ They continue to manage kindergartens
 - ▶ Only the headaches associated with the matching are taken away
- People do not like change
 - ▶ Start with a pilot in one or two councils

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Childcare Matching

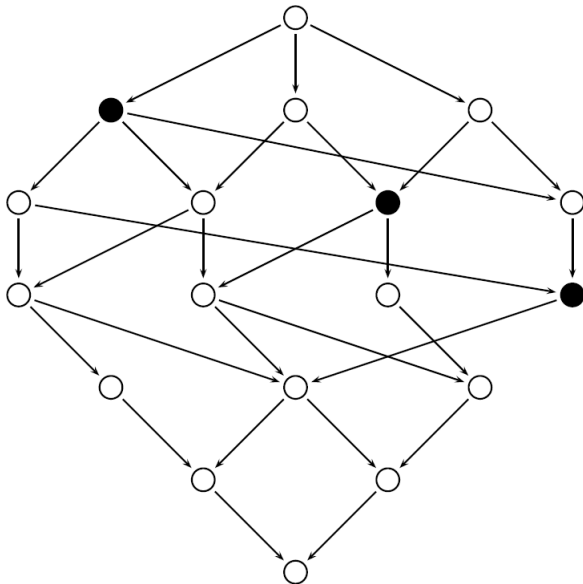
- A matching market
 - ▶ Families have preferences over childcare centres
 - ▶ Priorities are determined by law (centres can to some extent have a say)
 - ▶ Each centre has a capacity limit
- Two main differences with kindergarten
 - ▶ Children can enter or leave at any point
 - ▶ Children can attend part-time
- Dynamic issue
 - ▶ Trade-off in terms of how often the market is cleared
 - ▶ Thicker market vs waiting time
- Part-time issue
 - ▶ Enormous consequences on the model
 - ▶ This is what I study in my paper

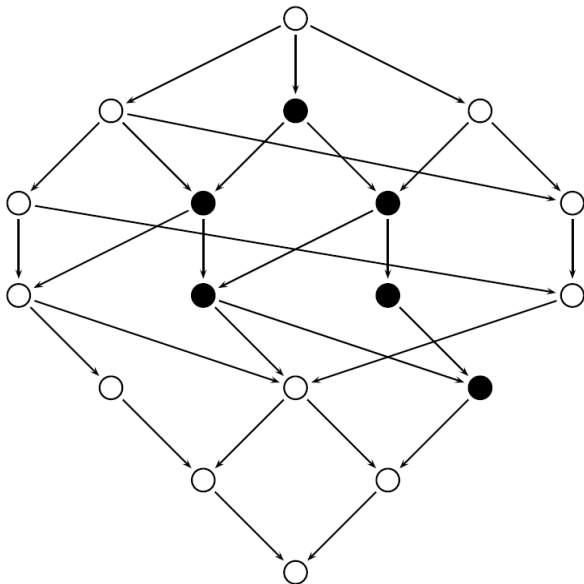
Matching with Quantity

- Simple model
 - ▶ Focuses on the heart of the problem
 - ▶ Any application, including childcare is inevitably more complex
 - ▶ The main insights developed are still valid
- Some agents want two units of the same good
 - ▶ These agents are not interested in getting just one unit
 - ▶ Children who need to attend childcare full-time
- Some agents only want one unit
 - ▶ Children who need to attend part-time
- Complementarity in preferences
 - ▶ An agent who wants two units sees them as complements
 - ▶ A unit is worth more to the agent if (s)he already has one
 - ▶ This is the heart of the problem

Consequences of Complementarity

- The set of stable matching is not well behaved
 - ▶ It may be empty
 - ▶ It may not contain an agent-optimal stable matching
 - ▶ Instead there may be several undominated ones
- The deferred acceptance algorithm does not work
 - ▶ Even if an agent-optimal stable matching exists it may not find it
- What is going on?
 - ▶ The set of stable matching is part of a larger set
 - ▶ That set is well behaved
- Relax the definition of stability
 - ▶ Allow for some degree infeasibility and instability
 - ▶ “Pseudo-stable” matchings
 - ▶ Search that well-behaved set for stable matchings





Finding Stable Matchings

- The algorithm works in two stages
 - (i) Adapt the deferred acceptance algorithm to find the agent-optimal pseudo-stable matching
 - (ii) Search the set to find stable matchings
- All stable matchings can be found in this way
 - ▶ This can be computationally heavy
 - ▶ Finding an undominated stable matching may be enough
- Applications
 - ▶ Childcare matching
 - ▶ University exchange programs
 - ▶ Matching with couples
 - ▶ **Refugee dispersal**

Refugee Dispersal

- Joint project
 - ▶ Scott Kominers (Harvard)
 - ▶ Alex Teytelboym (Oxford)
- The United Kingdom will resettle 20,000 Syrian refugees by 2020
 - ▶ These will be spread across the country in several localities
- We study this matching market
 - ▶ Refugees have preferences over localities
 - ▶ Localities can set up priorities
 - ▶ Localities have capacity limits

Refugee Dispersal

- Refugees are more likely to successfully integrate if
 - ▶ They are relocated in a place they like
 - ▶ They have the services they need
 - ▶ They have a chance to find work
 - ▶ They are a good fit for the community
- Technical Difficulty
 - ▶ Families have different sizes
 - ▶ Families require different services (schools, hospitals, etc)
- Complex version of the quantity problem
 - ▶ A similar algorithm can be found to find a stable matching

International Refugee Crisis

- Refugees currently have three options
 - ▶ Apply to one country at a time
 - ▶ Wait around in a camp to be processed by the UN
 - ▶ Reach Europe (or Australia) and claim asylum
- This could be organised as a matching market
 - ▶ Refugees have preferences over countries
 - ▶ Countries have preferences over refugees and set quotas
 - ▶ Quantity problem does not matter on such a large scale
- This is a standard two-sided matching market
 - ▶ Deferred acceptance works well
 - ▶ The hard part is to convince countries to offer resettlement places
 - ▶ The quotas must be high enough for refugees to enter the system rather than seek asylum

Conclusion

- Matching theory has many applications
 - ▶ School choice, kidney exchange, labour market, university admission, doctor-hospital matching, cadet matching, refugee dispersal, etc
 - ▶ Organising these markets efficiently can make a real difference
- Research continues
 - ▶ More complex matching models and algorithms are being developed
 - ▶ Potential for more applications
- The current theory already has great potential
 - ▶ It is underutilised and many markets could be improved
- Academics have little incentive to tackle these problems
 - ▶ This is the purpose of the Center for Market Design
 - ▶ Public servants have a very important role to play

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